

# Operating Systems 2013/14 0x4D4153 Assignment 4

Prof. Dr. Frank Bellosa Dipl.-Inform. Marius Hillenbrand Dipl.-Inform. Marc Rittinghaus

## Submission Deadline: Monday, January 13th, 2014 – 9:30 a.m.

In this assignment you will study sychronization, deadlocks, and memory management. All the organizational remarks of the first assignment are still valid!

Please print out the pages containing **T-Questions** and answer them on your printout. Clearly mark every page with your name, matriculation number and **tutorial number**. Simply put it in the mailbox in the basement of building 50.34 (Info-Neubau).

**P-Guestions** are programming assignments. Download the provided tarball from the VAB and make sure to use the included templates and Makefiles. Do not fiddle with the compiler flags. Submission instructions can be found in the first assignment.

#### Any assignment handed in after its deadline will be ignored!

## **T-Question 4.1: Deadlocks**

a. Analyze the code fragment below. Can a deadlock occur? Why, or why not?

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```
Mutex m1, m2, m3 = 1;
Thread1()
                                            Thread2()
{
                                            ł
    wait(m1):
                                                wait (m2);
    // update some data
                                                // update some data
    signal(ml);
                                                wait (m3);
    wait (m2);
                                                // update some more data
    wait (m3):
                                                signal(m2);
    // update some more data
                                                wait(m1):
    signal(m2);
                                                // update even more data
    signal(m3);
                                                signal(m3);
}
                                                signal(m1);
                                            }
```

b. Explain why an unsafe state does not always lead to a deadlock!

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c.	Given a system with 3 processes, $P_1$ to $P_3$ , and three resource types $R_1$ to $R_3$ . Assume that there is only one instance of each type. Depict the resource-allocation graph for the following situation: $P_1$ has requested $R_1$ and $R_2$ . $P_2$ is holding $R_1$ and is waiting for $R_2$ . $P_3$ has acquired $R_3$ and is waiting for $R_3$ .	1 T-nt
	is waiting for <i>1</i> (3, <i>1</i> 3 has acquired <i>1</i> (3 and 15 waiting for <i>1</i> (2).	I I pt
d	Has a deadlook accurred in the above situation? Why, or why not? (Assume that	
u.	resources can neither be shared nor preempted.)	1 T-pt
e	The processes $P_1$ , $P_2$ , and $P_3$ described above only need the resources they have	
c.	already requested or acquired already. Is the system in a safe state? Prove your claim!	9 T-nt
		2 1-pt
т	Organtian A.D. Manager Allagation	
1	What HW support is required for a simple partition based memory management?	1 T at
a.	what new support is required for a simple partition-based memory management?	1 1-pt
b.	What is the difference between internal and external fragmentation?	1 T-pt

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c.	Why would you prefer the allocation policy first fit over best fit?	1 T-pt
d.	Why might best fit produce more fragmentation than first fit in some scenarios?	1 T-pt
<b>T</b> -	<b>Question 4.3: Segmentation</b>	1 T-nt
a.		<b>F</b>
b.	Consider a system with 16-bit logical addresses that supports 4 different segments. What is the maximum size a segment can have?	1 T-pt

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c. Assume a system with 16-bit logical addresses that supports four different segments, which uses the following segment table:

segment no.	base	limit
0	Oxdead	0xbeef
1	0xf154	0x013a
2	0x0000	0x0000
3	0x0000	0x4711

Complete the following table and explain briefly how you derived your solution for each row in the table.

logical address	segment number	offset	valid?	physical address
	3	0x3999		
0x2020				
0x9859				
			yes	0xf15f



## **T-Question 4.4: Bonus Questions**

a.	A spinlock t	hat is frozen in	place must b	e nuked from orb	oit. Explain	why.	0 T-pt
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b. How does an OS handle the diversity in the gap? (Note: There is no orbit around the gap.)

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## **P-Question 4.1: Contiguous Allocation**

You are in a team that implements an OS that uses contiguous allocation for process memory. Your task is to write the allocator that allocates and frees memory partitions of arbitrary sizes (with a granularity of 1 kB). You decide to develop and test your allocator as a user-level process before integrating it into your new OS.

Your allocator will manage memory of size MEM\_SIZE kB that you first allocate on the heap using malloc. Your allocator should label each memory partition with the identifier of the process running in that partition. For easier testing, you chose to identify processes with letters (A, B and so on). Make sure that your code works independently of the actual value of the macro MEM\_SIZE.

Add your functions to the file allocator.c, but keep the function signatures in allocator.h unchanged. You may add test code to main.c.

a. First, design and implement the data structures you want to use in your memory allocator. Then, add a function that prints all current allocations in a *memory map* to help you debug your code: For each kB in the managed memory, print the partition it is assigned to, ordered from address 0 to MEM\_SIZE-1 kB. Print unallocated space as '0'. For example, three partitions 'A', 'B', and 'C', with sizes of 3, 1, and 2, might occur in the memory map as AAAOCCB000...0.

b.	Implement the <i>first fit</i> allocation policy and place it in <code>allocate_partition_first_fit</code> . Test your solution using the memory map function. Try to provoke fragmentation.	2 P-pt
c.	Your project management decides to use the <i>best fit</i> policy instead. Implement the new policy in the function allocate_partition_best_fit.	2 P-pt
d.	Finally, also implement the <i>worst fit</i> allocation policy as a third alternative in func- tion allocate_partition_worst_fit.	2 P-pt

## **P-Question 4.2: Segmentation**

Implement the address translation mechanism of segmentation. Follow the scheme in T-Question 4.3: Use 16-bit logical addresses and four different segments.

Put your solution into segment.c. Do not change segment.h. The segment table is filled in main.c where you may add arbitrary test code (try the values from the theory question). However, your solution must work with any values in the segment table.

- a. Write code that calculates the following values from a logical address and a segmentation table:
  - the segment number,
  - the offset,
  - the validity of the logical address, and
  - the physical address (or 0 for invalid logical addresses)

Fill out the function stubs in segment.c.

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#### **P-Question 4.3: Shared Memory**

Have a look at asst4-sharedmem/sharedmem.c. This code template forks a child process that is intended to act as a client, while the parent process represents the server. Your job is to implement both client and server and the communication between the two, using POSIX shared memory.

Use the lecture slides on process coordination as a starting point to find out how to use shared memory. Also, have a look at the man-pages of shmget, shmat, and shmdt.

- a. Establish a shared memory segment before forking the client so that both server and client know the id of the shared segment. For every shared memory call you do, write a comment, that explains what the passed arguments do and why you pass them! You will not be accredited points for the calls that do not have this mandatory comment.
- b. Implement the client() function to read two unsigned integer values from standard input and pass them to the server via the shared memory segment.Wrap all values you pass through the shared memory segment in one single struct! The shared memory segment may only contain a single variable (the struct) for a legal solution!
- c. Implement the server() function to read the two values written by the client, calculate the product of the two values, and print the result followed by a newline to standard output. Limit the output of your solution to the mentioned elements. Don't print more!
- d. Use POSIX semaphores for signaling between the client and the server process: The server must wait until the client entered two values before processing the input, the client must wait until the server is done processing the input and printing the output before prompting for new input values. Hint: Use two semaphores.
- e. Entering -1 for both values shall terminate both the client and the server process. Take care of cleaning up the shared memory segment.

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